

Nim - the first high performance language with full support for hot code-reloading at runtime

by Viktor Kirilov

Me, myself and I

- my name is Viktor Kirilov - from Bulgaria
- creator of `doctest` - the **fastest** C++ testing framework
- apparently I like text-heavy slides and reading from them...!
 - deal with it :|

Talk agenda

- some Nim code
- the performant programming language landscape
 - read: heavily biased C++ rant
- Nim compilation model
- hot code reloading
 - usage & implementation
 - ".dll" => assume .so/.dylib (platform-agnostic)
- demo
- comments & conclusions
- a bit on REPLs

Hello

```
1 echo "Hello World"
```

Currencies

```
1 type
2   # or use {.borrow.} here to inherit everything
3   Dollars* = distinct float
4
5 proc `+` *(a, b: Dollars): Dollars {.borrow.}
6
7 var a = 20.Dollars
8
9 a = 25 # Doesn't compile
10 a = 25.Dollars # Works fine
11
12 a = 20.Dollars * 20.Dollars # Doesn't compile
13 a = 20.Dollars + 20.Dollars # Works fine
```

Sets

Operator	Description	Example Code
<code>a in B</code>	is a an element of B?	<code>'d' in {'a'..'z'}</code>
<code>a notin B</code>	is a not an element of B?	<code>40 notin {2..20}</code>
<code>A + B</code>	union of A with B	<code>{'a'..'m'} + {'n'..'z'} == {'a'..'z'}</code>
<code>A - B</code>	relative complement of A in B	<code>{'a'..'z'} - {'b'..'d'} == {'a', 'e'..'z'}</code>
<code>A + {b}</code>	add element b to set A	<code>{'b'..'z'} + {'a'} == {'a'..'z'}</code>
<code>A - {b}</code>	remove element b from set A	<code>{'a'..'z'} - {'a'} == {'b'..'z'}</code>
<code>A * B</code>	intersection of A with B	<code>{'a'..'m'} * {'c'..'z'} == {'c'..'m'}</code>
<code>A <= B</code>	is A a subset of B?	<code>{'a'..'c'} <= {'a'..'z'}</code>
<code>A < B</code>	is A a strict subset of B?	<code>{'b'..'c'} < {'a'..'z'}</code>

Iterators

```
1 type
2     CustomRange = object
3         low: int
4         high: int
5
6 iterator items(range: CustomRange): int =
7     var i = range.low
8     while i <= range.high:
9         yield i
10        inc i
11
12 iterator pairs(range: CustomRange): tuple[a: int, b: char] =
13     for i in range: # uses CustomRange.items
14         yield (i, char(i + ord('a')))
15
16 for i, c in CustomRange(low: 1, high: 3):
17     echo c
18
19 # prints: b, c, d
```

Variants

```
1 # This is an example how an abstract syntax tree could be modelled in Nim
2 type
3   NodeKind = enum # the different node types
4     nkInt,      # a leaf with an integer value
5     nkFloat,   # a leaf with a float value
6     nkString,  # a leaf with a string value
7     nkAdd,     # an addition
8     nkSub,     # a subtraction
9     nkIf       # an if statement
10  Node = ref object
11    case kind: NodeKind # the ``kind`` field is the discriminator
12    of nkInt: intVal: int
13    of nkFloat: floatVal: float
14    of nkString: strVal: string
15    of nkAdd, nkSub:
16      leftOp, rightOp: Node
17    of nkIf:
18      condition, thenPart, elsePart: Node
19
20  var n = Node(kind: nkFloat, floatVal: 1.0)
21  # the following statement raises an `FieldError` exception, because
22  # n.kind's value does not fit:
23  n.strVal = ""
```

Multi methods

```
1 type
2   Thing = ref object of RootObj
3   Unit = ref object of Thing
4     x: int
5
6 method collide(a, b: Thing) {.inline.} =
7   quit "to override!"
8
9 method collide(a: Thing, b: Unit) {.inline.} =
10  echo "1"
11
12 method collide(a: Unit, b: Thing) {.inline.} =
13  echo "2"
14
15 var a, b: Unit
16 new a
17 new b
18 collide(a, b) # output: 2
```

Meta-programming

- what is it
 - a program that can read, generate, analyze or transform other programs
- why do it
 - can optimise code – by compile-time rewrites
 - think expression templates
 - can enforce better coding patterns
 - can increase code readability and maintainability
 - with great power comes great responsibility
- reflection - when the meta language is the actual language

Meta-programming in Nim

- works on the Abstract Syntax Tree
- respects the type system
- levels of complexity:
 - normal procs and inline iterators
 - generic procs and closure iterators
 - templates
 - macros

Templates

```
1 template withFile(f: untyped, filename: string,  
2                 mode: FileMode,  
3                 body: untyped): typed =  
4   let fn = filename  
5   var f: File  
6   if open(f, fn, mode):  
7     try:  
8       body  
9     finally:  
10      close(f)  
11  else:  
12    quit("cannot open: " & fn)  
13  
14 withFile(txt, "ttempl3.txt", fmWrite):  
15   txt.writeLine("line 1")  
16   txt.writeLine("line 2")
```

AST

```
1 dumpTree:
2   var mt: MyType = MyType(a:123.456, b:"abcdef")
3
4 # output:
5 #   StmtList
6 #     VarSection
7 #       IdentDefs
8 #         Ident "mt"
9 #         Ident "MyType"
10 #       ObjConstr
11 #         Ident "MyType"
12 #         ExprColonExpr
13 #           Ident "a"
14 #           FloatLit 123.456
15 #         ExprColonExpr
16 #           Ident "b"
17 #           StrLit "abcdef"
```

Macros

```
1 import macros
2
3 type
4   MyType = object
5     a: float
6     b: string
7
8 macro myMacro(arg: untyped): untyped =
9   var mt: MyType = MyType(a:123.456, b:"abcdef")
10  let mtLit = newLit(mt)
11
12  result = quote do:
13    echo `arg`
14    echo `mtLit`
15
16 myMacro("Hallo")
17
18 # The call to myMacro will generate the following code:
19 echo "Hallo"
20 echo MyType(a: 123.456'f64, b: "abcdef")
```

More macros

```
1 import macros
2 dumpAstGen:
3   proc hello() =
4     echo "hi"
```

```
1 nnkStmtList.newTree(
2   nnkProcDef.newTree(
3     newIdentNode(!"hello"),
4     newEmptyNode(),
5     newEmptyNode(),
6     nnkFormalParams.newTree(
7       newEmptyNode()
8     ),
9     newEmptyNode(),
10    newEmptyNode(),
11    nnkStmtList.newTree(
12      nnkCommand.newTree(
13        newIdentNode(!"echo"),
14        newLit("hi")
15      )
16    )
17  )
18 )
```

More macros - continue from last slide

```
1 import macros
2 macro gen_hello(): typed =
3     result = nnkStmtList.newTree(
4         nnkProcDef.newTree(
5             newIdentNode(!"hello"),
6             newEmptyNode(),
7             newEmptyNode(),
8             nnkFormalParams.newTree(
9                 newEmptyNode()
10            ),
11            newEmptyNode(),
12            newEmptyNode(),
13            nnkStmtList.newTree(
14                nnkCommand.newTree(
15                    newIdentNode(!"echo"),
16                    newLit("hi")
17                )
18            )
19        )
20    )
21 gen_hello()
22 hello() # << same as from last slide!
```

HTML DSL

```
1 import html_dsl
2
3 html page:
4   head:
5     title("Title")
6   body:
7     p("Hello")
8     p("World")
9     dv:
10      p "Example"
11
12 echo render(page())
```

HTML DSL result

```
1 <!DOCTYPE html>
2 <html class='has-navbar-fixed-top' >
3 <head>
4   <meta charset="utf-8">
5   <meta name="viewport" content="width=device-width, initial-scale=1">
6   <title>Title</title>
7 </head>
8 <body class='has-navbar-fixed-top' >
9   <p >Hello</p>
10  <p >World</p>
11  <div>
12    <p>Example</p>
13  </div>
14 </body>
15 </html>
```

Simply Nim

- statically typed
- high performance (compiles to native binaries - comparable to C/C++)
- very clean & elegant - no, beauty is NOT subjective!
- garbage collected (can do manual memory management too)
- expressive - some of the most powerful metaprogramming
 - compiler has an interpreter inside
- compiles to: C/C++/ObjC/Javascript
 - non-idiomatic - not for reading but optimal for execution
- suited for: systems programming, applications & web
 - all types of software!
- backed by Status since 2018 (#65 cryptocurrency by marketshare)
 - Status - working on one of the first implementations of Ethereum 2.0
 - just like Rust is backed by Mozilla (although with a lot less...)
- has a rich stdlib, package manager, docs, some IDE support

Feature rundown

- uniform call syntax (extension methods) - `obj.method()` OR `method(obj)`
 - that's why there are no real "methods" defined in types
- function call parens are optional - `echo("hello")` OR `echo "hello"`
- case-insensitive - also underscore-insensitive but that's another topic :|
- generics
- templates (meta-programming^{^2})
- macros (meta-programming^{^3}) - evaluated in the compiler by the NimVM
- concepts
- discriminated unions
- strong typedefs (distinct type) - can has \$ currency?
- coroutines & closures
- switch & pattern matching
- dynamic dispatch & multi-methods
- converters - explicit (for implicit conversions)
- effect system (transitive)
- extensible pragmas, "defer", exceptions, "discard", named args... good defaults!

My "favourite" aspect of C++

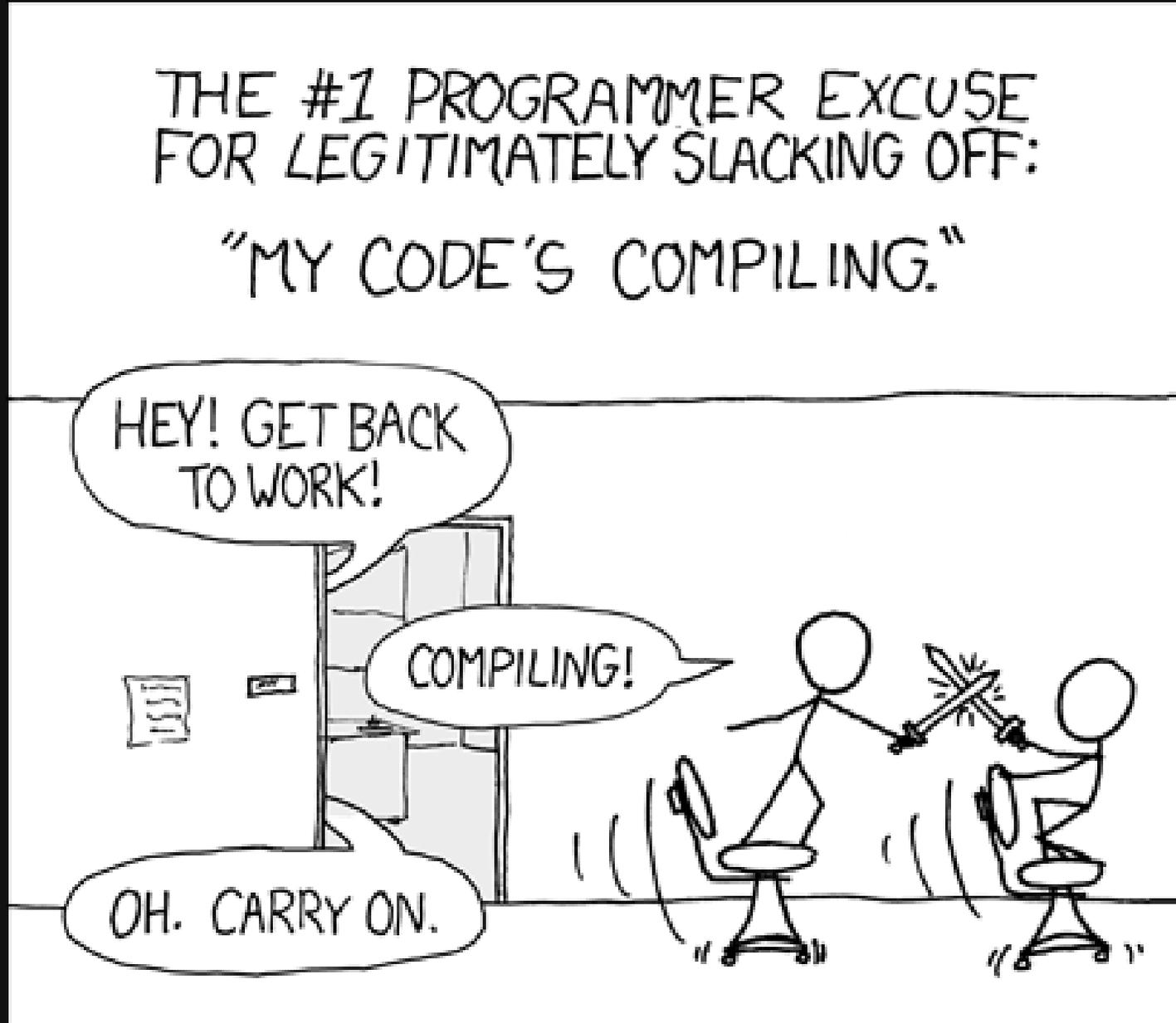
THE #1 PROGRAMMER EXCUSE
FOR LEGITIMATELY SLACKING OFF:

"MY CODE'S COMPILING."

HEY! GET BACK
TO WORK!

COMPILING!

OH. CARRY ON.



A bit on C++

- C++20 is shaping up to be a huge release
 - lots of cool stuff, but complexity is through the roof
 - Expert-"tolerable" - prestige when you come up with yet more complicated TMP
- simple example using ranges from C++20 - [blog post](#)
 - 3 seconds of compile time for ~20 lines of code, forget about "Debug" builds
- [Remember the Vasa!](#) - Bjarne Stroustrup
- There should come a time for a clean slate
 - C++ is a great and valuable ongoing research
 - The 2 biggest reasons C++ is so widely used today:
 - legacy and maturity - too much software written already
 - inertia - attachment and lack of interest to learn new languages
 - C++ is a HUGE time/money cost on the scale of hundreds of millions
 - developer productivity, bug & safety
 - business should back a better language & push for development + learning

Some quotes & thoughts

- Fifty years of programming language research, and we end up with C++?
 - Richard A. O'Keefe
- There are only two kinds of programming languages: those people always bitch about and those nobody uses.
 - Bjarne Stroustrup
- Nim is the next iteration of practical language design
 - by humble !!! >> me << !!!
- Nim: speed of C, elegance of Python, flexibility of Perl
 - Peter Munch-Ellingsen
- Nim is to C++ as CoffeeScript is to JavaScript
 - [cjhanks, hackernews](#) Apr 18, 2017

Comparison with others

- D, Rust, Jai, Zig
 - out of scope for this talk
- Go
 - not really a *pinnacle* of abstraction and innovation :|
- C++
 - <optional> - 5k+ LOC for a T and a bool... safe_int - same horror story
 - The next big thing: "Design by introspection" - Andrei Alexandrescu
- Nim is one of the most logical paths forward
 - on-par performance with C/C++ (compiles to them)
 - some of the most easy interop with C/C++ (compiles to them)
 - uses any C/C++ compiler (compiles to them)
 - already quite far in terms of implementation
 - meta-programming on steroids

Nim compilation model

```
1 # main.nim
2
3 import a
4
5 echo a()
```

```
1 # a.nim
2
3 import b
4
5 proc a*(): string =
6   return from_b
```

```
1 # b.nim
2
3 let local = "B!"
4
5 let from_b* = local
6   ^
7 # means "exported"
```

- `nim c -d:release main.nim`
 - always compile only the main file, follow the imports
 - whole program analysis
 - a `.c` file for each `.nim` file in a "nimcache" (temp) folder (also `.obj` files)
 - only referenced (imported) modules are compiled in the end
- entire project is always "compiled" by Nim (currently no "minimal" rebuild)
 - ~4-5 sec for the entire source of Nim - 135 files (without the C compiler)
 - the C/C++ compiler rebuilds only changed files (takes a bit more time)
 - will change when per-module caching is introduced - even faster!

Nim to C/C++: nimbase.h

included by all .c/.cpp files

```
1 // nimbase.h
2
3 #define N_NIMCALL(rettype, name) rettype __fastcall name
4 #define N_CDECL(rettype, name) rettype __cdecl name
5 //...
6 #define N_NIMCALL_PTR(rettype, name) rettype (__fastcall *name)
7 //...
8 #define N_LIB_PRIVATE __attribute__((visibility("hidden")))
9 //...
10 #define N_LIB_EXPORT extern __declspec(dllexport)
11 //...
12 #define STRING_LITERAL(name, str, length) \
13     static const struct { \
14         TGenericSeq Sup; \
15         char data[(length) + 1]; \
16     } name = {{length, (int) ((unsigned)length | NIM_STRLIT_FLAG)}, str}
```

handles different platforms - convenience macros

Nim procs to C/C++

```
1 proc foo() =
2   echo "hello"
3
4 foo()
```

```
1 #include <nimbase.h>
2
3 // forward declarations / type definitions / constants section
4 struct TGenericSeq { int len; int reserved; };
5 struct NimStringDesc : public TGenericSeq { ... };
6 typedef NimStringDesc* tyArray_nHXaesL0DJZHyVS07ARPRA[1];
7
8 STRING_LITERAL(TM_r9bkcJ6PRJ5n7ORNxxJ5ryg_3, "hello", 5); // << string literal
9 NIM_CONST tyArray_nHXaesL0DJZHyVS07ARPRA TM_r9bkcJ6PRJ5n7ORNxxJ5ryg_2 =
10   {((NimStringDesc*) &TM_r9bkcJ6PRJ5n7ORNxxJ5ryg_3)};
11
12 N_LIB_PRIVATE N_NIMCALL(void, foo_iineYNh8S9cE6Ry7dr2Tz2A)(void); // << fwd decl
13
14 // definition section
15 N_LIB_PRIVATE N_NIMCALL(void, foo_iineYNh8S9cE6Ry7dr2Tz2A)(void) { // << def
16   echoBinSafe(TM_r9bkcJ6PRJ5n7ORNxxJ5ryg_2, 1); // the echo call
17 }
18
19 // code execution section
20 foo_iineYNh8S9cE6Ry7dr2Tz2A(); // << call
```

Nim types to C/C++

```
1 type
2   MyData = object
3     answer: int
4     ready: bool
5 proc newData(): MyData = return MyData(answer: 42, ready: true)
6 echo newData().answer
```

```
1 // forward declarations / type definitions / constants section
2 struct tyObject_MyData {
3     int answer;
4     bool ready;
5 };
6 // definition section
7 N_LIB_PRIVATE N_NIMCALL(tyObject_MyData, newData)(void) {
8     tyObject_MyData result; // always an implicit "result"
9     nimZeroMem((void*)&result, sizeof(tyObject_MyData));
10    result.answer = ((int) 42);
11    result.ready = true;
12    return result;
13 }
14
15 // code execution section
16 tyObject_MyData T2_;
17 T2_ = newData(); // << call
18 //...
```

Nim closures to C/C++ (resumable funcs)

```
1 iterator closure_iter*(): int {.closure.} = # a resumable function
2   var x = 1
3   while x < 10:
4     yield x
5     inc x
6 for i in closure_iter(): echo i
```

```
1 struct state_type : public RootObj {
2   int colonstate_; // state progress - there are some GOTOs using this
3   int x1; // the state
4 };
5
6 struct closure_type {
7   N_NIMCALL_PTR(int, c_ptr) (void* e_ptr); // function ptr
8   void* e_ptr; // environment ptr
9 };
10
11 N_LIB_PRIVATE N_CLOSURE(int, func)(void* e_ptr) { // def omitted for simplicity
12
13 state_type st; // the state
14 closure_type local; // the closure
15 local.c_ptr = func; // assign the func
16 local.e_ptr = &st; // assign environment
17 //...
18 i = local.c_ptr(local.e_ptr); // the call in the loop
```

Nim compilation to C/C++: a BIG win

- smaller scope for the compiler
- all the cutting-edge optimization for C/C++ for free
- out-of-the-box support for tons of platforms
- easiest C/C++ interop possible
- exceptions - reusing those of C++ when using that backend
- nim to C/C++ code mapping with #line directives for debuggers
- no generated headers for the exported parts of modules
- each .c/.cpp file contains everything (and only what) it needs
 - forward declarations for external functions
 - type definitions
- each .c/.cpp file includes nimbase.h and a few C stdlib headers
- high level macros & templates => simple structs and functions

Interfacing with C/C++

Foreign Function Interface

```
proc printf(formatstr: cstring)
  {.header: "<stdio.h>", importc: "printf", varargs.}
```

other pragmas - for use in Nim:

```
{.emit: ""
using namespace core;
"".}

{.compile: "logic.c".}
```

We can also call Nim code from C/C++:

```
# fib.nim
proc fib(a: cint): cint {.exportc.} # do not mangle
```

```
nim c --noMain --noLinking --header:fib.h fib.nim
```

```
// user.c
#include <fib.h>
```

Interfacing with C/C++

C++ template constructs

```
1 type
2   StdMap {.importcpp: "std::map", header: "<map>".} [K, V] = object
3 proc `[ ]=`[K, V](this: var StdMap[K, V]; key: K; val: V) {.
4   importcpp: "#[#] = #", header: "<map>".}
5
6 var x: StdMap[cint, cdouble]
7 x[6] = 91.4
```

Generated C++

```
1 std::map<int, double> x;
2 x[6] = 91.4;
```

c2nim tool - generate C/C++ bindings for Nim

Runtime compilation - WHY

- much faster iteration times
 - no need to restart the program - can preserve state
- less need for a scripting language
 - no need for a virtual machine
 - no binding layer
 - code in one language
- can hack something quickly
 - introspection, queries
 - debuggers aren't infinitely powerful
 - fine-tuning values
- interactive (REPL-like): very useful for exploration and teaching

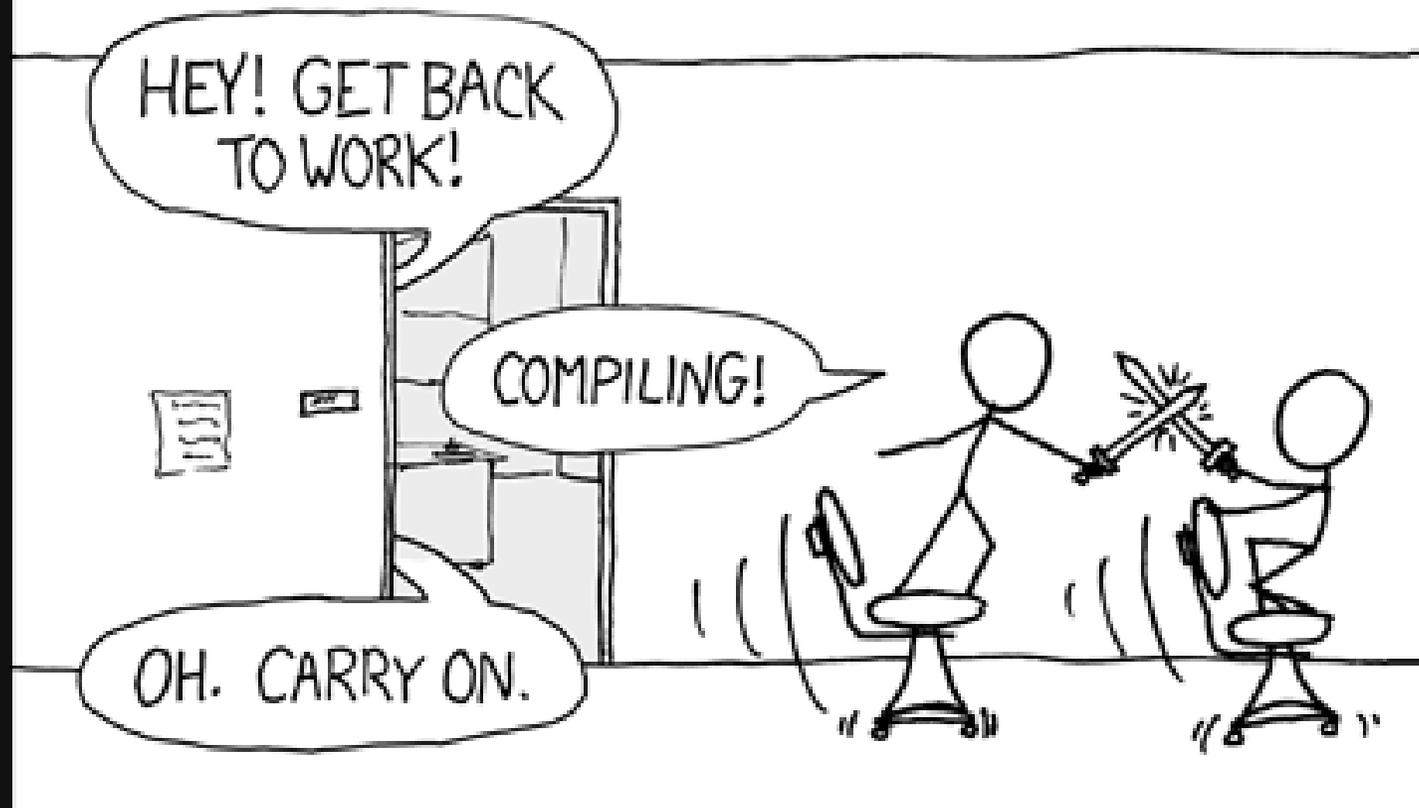
Runtime compilation for C/C++: HOW

- replacing entire functions: using shared libraries OR hot-patching:
 - possible for decades - but not widely used
 - usually quite intrusive (interfaces, constraints, complicated setup)
 - in game engines: Unreal, others...
 - hot-patching (with very little setup): [Live++](#), [Recode](#)
 - Visual Studio "Edit & Continue" - 0 setup, but limited
 - <https://github.com/crosire/blink>
 - <https://github.com/ddovod/jet-live>
 - <http://bit.ly/runtime-compilation-alternatives> << "one link to rule them all"
- interactive: REPL-like
 - [cling](#) - by researchers at CERN - built on top of LLVM
 - [inspector](#), [Jupiter](#)
 - hard to integrate in a platform/compiler agnostic way
 - [RCRL](#) - basically a hack - the inspiration for the Nim implementation

Replace "compiling" with "restarting"

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Hot code-reloading (HCR) in Nim

- inspired by a hacky REPL for C++ (called RCRL - by me)
- <https://github.com/nim-lang/Nim/issues/8927>
 - mentored by Zahary
- compile with `--hotCodeReloading:on`
- need also 2 .dlls (the HCR runtime + the GC of Nim)

```
1 # main.nim
2
3 import hotcodereloading # for reload
4 import other
5
6 while true:
7   echo readLine(stdin) # pause
8   performCodeReload() # reload
9   echo getInt()       # call
```

built as an .exe/.dll depending
on the project type

```
1 # other.nim
2
3 import hotcodereloading # for after handler
4
5 var glob = 42
6
7 proc getInt*(): int = return glob + 1 # exported
8
9 afterCodeReload:
10   glob = 666
```

built as a reloadable .dll
ends up in the "nimcache"

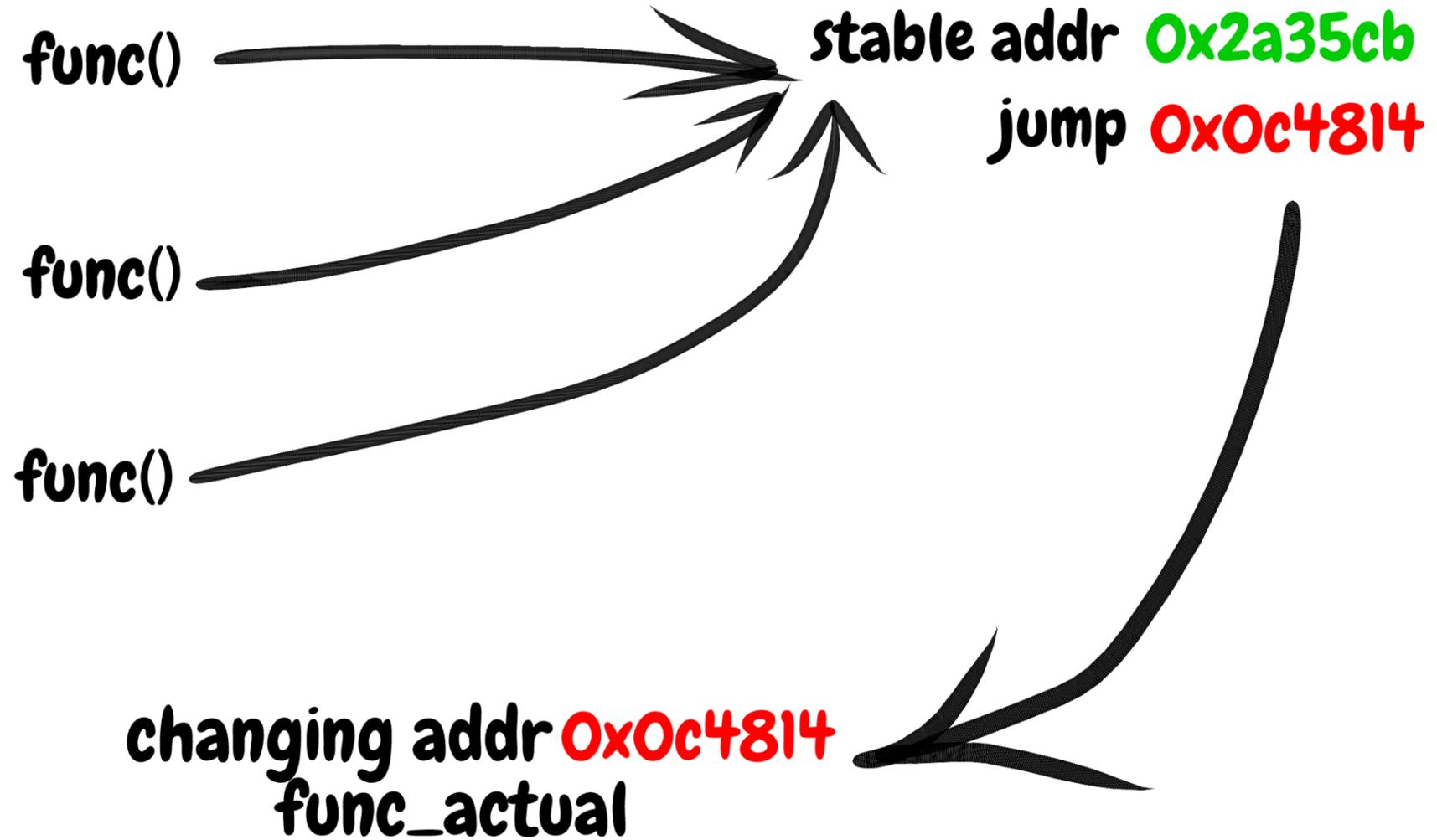
Effects of HCR

- all interaction between .nim modules => through pointers
- functions - changes:
 - forward declarations become function pointers
 - definitions get "_actual" as a suffix
 - pointers are assigned the "_actual" on startup
 - calls stay the same (pointer has the same name)
- globals - changes:
 - turned into pointers
 - allocated on the heap and initialized on startup
 - state is preserved when reloading
 - dereferenced wherever used

Effects of HCR

```
1 // fwd decl/globals section
2 static N_NIMCALL_PTR(int, getInt_omy6T2FkprLEReOy2ITmIQ)(void);
3 static int* glob_v1zK9aUOu9aNNcsxruuK8NdA;
4
5 // definitions
6 N_LIB_PRIVATE N_NIMCALL(int, getInt_omy6T2FkprLEReOy2ITmIQ_actual)(void) {
7     int result; // ^^ the suffix
8     result = (*glob_v1zK9aUOu9aNNcsxruuK8NdA);
9     return result;
10 }
11
12 // usage
13 (*glob_v1zK9aUOu9aNNcsxruuK8NdA) = getInt_omy6T2FkprLEReOy2ITmIQ();
14
15 // init on startup (naive)
16 glob_v1zK9aUOu9aNNcsxruuK8NdA = new int(42);
17 getInt_omy6T2FkprLEReOy2ITmIQ = getInt_omy6T2FkprLEReOy2ITmIQ_actual
```

Trampolines



Initialization

```
1 // naive
2 glob_v1zK9aUOu9aNNcsxruuK8NdA = new int(42);
3 getInt_omy6T2FkprLEReOy2ITmIQ = getInt_omy6T2FkprLEReOy2ITmIQ_actual
```

```
1 // reality
2 getInt_omy6T2FkprLEReOy2ITmIQ = (tyProc_vVu2P82aVLv9c8X0xbI1NJw) hcrRegisterProc(
3     "D:\\play\\nimcache/play.cpp.dll",           // "domain" (AKA module)
4     "getInt_omy6T2FkprLEReOy2ITmIQ",           // "key"
5     (void*)getInt_omy6T2FkprLEReOy2ITmIQ_actual); // the real function
6
7 if(hcrRegisterGlobal("D:\\play\\nimcache/play.cpp.dll",           // "domain" (AKA module)
8     "glob_v1zK9aUOu9aNNcsxruuK8NdA",           // "key"
9     sizeof((*glob_v1zK9aUOu9aNNcsxruuK8NdA)), // size for allocation
10    NULL, // for the GC - simple integer is simple, so NULL
11    (void**)&glob_v1zK9aUOu9aNNcsxruuK8NdA) // address to pointer
12 {
13     // hcrRegisterGlobal returns "true" only if not already initied
14     (*glob_v1zK9aUOu9aNNcsxruuK8NdA) = ((int) 42); // init with value (or side effects)
15 }
```

Initialization

- the HCR.dll runtime holds pointers to all globals/functions
- hcrRegisterProc
 - allocates executable memory (a few bytes)
 - writes a jump instruction (trampoline) to the "_actual"
 - returns an address to the trampoline
 - this way "_actual" can be changed on reloading
 - changed by calling it again with a different address
 - all pointers to the trampoline stay the same
- all symbols are registered per "domain" (.dll)
 - no name clashes (even though they are mangled...)
 - better management - can remove all symbols for module X

Initialization

```
1 # main.nim
2
3 import a, b
4
5 echo from_a()
6 echo from_b()
```

```
1 # a.nim
2
3 import b
4
5 proc from_a*(): string =
6   result = "A!"
7   result.add from_b()
```

```
1 # b.nim
2
3 proc from_b*(): string =
4   return "B!"
```

1. `main.exe` loads the `hcr.dll` (and the Nim GC in `rtl.dll`)
2. `main.exe` calls `init()` from `hcr.dll` and passes a list of imports (`a, b`)
3. `hcr.dll` loads `a.dll` and gets a list of imports (`b`)
4. `hcr.dll` loads `b.dll` and fully initializes it (it has no imports)
 1. registers `from_b()` and does nothing else
5. `hcr.dll` fully initializes `a.dll`
 1. registers `from_a()` and gets the address for `from_b()`
6. `hcr.dll` skips `b.dll` (part of the imports of `main.exe`) since it is already initialized
7. `main.exe` is initialized
 1. gets the addresses for `from_a()` and `from_b()`
 2. executes the top-level code (the 2 `echo` statements)

Initialization

- a DFS traversal with POST visit
- when module A imports a symbol from B
 - symbol is first registered in B
 - symbol is "gotten" in A after B is initied
- basically a custom dynamic linker :|
- imports are discovered on-the-go
- HCR.dll constructs a tree of imports and maintains it
- many details omitted
 - initialization is broken into multiple passes
 - registration of type infos (for the GC) is a pre-pass
- each .dll exports just a few functions which the HCR.dll uses
 - getImports(), and the ones for the passes

Reloading

when we call `performCodeReload()`:

- HCR.dll will check `hasAnyModuleChanged()`
 - basically scanning if any .dll has been modified (timestamp)
- changes shouldn't affect .dll files which are part of the current active callstack when `reload()` is called! or crash :|
 - ==> main module can never be reloaded
- execute the "beforeCodeReload" handlers if about to reload
- in a DFS traversal, for each modified module:
 - same as the init - get its imports, load them (if changed or new), init everything in proper order
 - supports discovery of new imports!
 - also removes no longer referenced modules and their symbols
- execute the "afterCodeReload" handlers

Reloading - handlers

```
1 # main.nim
2 import a, b, hotcodere reloading
3
4 beforeCodeReload:
5   echo "before main"
6 afterCodeReload:
7   echo "after main"
```

```
1 # a.nim
2 import b, hotcodere reloading
3
4 beforeCodeReload:
5   echo "before a"
6 afterCodeReload:
7   echo "after a"
```

```
1 # b.nim
2 import hotcodere reloading
3
4 beforeCodeReload:
5   echo "before b"
6 afterCodeReload:
7   echo "after b"
```

- DFS traversal with POST visit
- handlers can be added/removed
- can be used to update globals
- fine-grained control:
 - `hasModuleChanged(<module>)`

only A changes => all handlers
are executed on reload:

```
before b
before a
before main
after b
after a
after main
```

Reloading - global scope

- top-level code (global scope) is executed only on initial load
 - for new top-level code use before/after handlers
- changing the initializer of a global doesn't do anything
 - use a before/after handler
 - or remove the global entirely, reload, and re-add it
 - brand new symbol!
- new globals can be added - and will be initialized properly

The initial HCR example revisited

```
1 # main.nim
2
3 import hotcodereloading # for reload
4 import other
5
6 while true:
7   echo readLine(stdin) # pause
8   performCodeReload() # reload
9   echo getInt() # call
```

```
1 # other.nim
2
3 import hotcodereloading # for after handler
4
5 var glob = 42
6
7 proc getInt*(): int = return glob # exported
8
9 afterCodeReload:
10   glob = 666
```

Makes more sense now, doesn't it?

A dramatic space scene featuring an astronaut in a white suit floating in the void. The Earth's blue and white horizon is visible on the left, with a bright sun or light source creating a lens flare effect. Debris and particles are scattered throughout the scene, suggesting a high-speed or chaotic environment. The text 'LIVE DEMO' is centered in the image.

**LIVE
DEMO**

Encountered problems

- processes lock loaded .dll files in the filesystem on Windows
 - when reloading we copy x.dll to x_copy.dll and load the copy
- changing module X can affect module Y
 - such changes shouldn't reach the main module
 - mangling of symbols being affected by attributes (purity)
 - mangling affected by where "inline" functions get used first
 - mangling affected by which module instantiates a generic
- C vs C++
 - missing forward declarations - fine in C!
 - multiple identical forward declarations
 - multiple definitions of global function pointers - fine in C!

Visual Studio debug symbols - PDB drama

- .dll/.exe have hardcoded paths to the .pdb (copying the .dll doesn't matter)
- the VS Debugger keeps the .pdb files locked for .dlls even after unloaded

solutions:

- someone managed to **close the file handles** to no longer needed .pdb files (.dll has been unloaded) to the external VS debugger process (**live++**)
- embed the debug info in the actual binaries just like on unix
 - /Z7 embeds it in .obj files but not for the final .dll/.exe when linking them
- different names for the .pdb using /PDB:<filename> (with the date/time (including milliseconds) as a suffix)
 - the "hardcoded" paths to .pdb files are always different
 - try to delete all <dll_name>_*.pdb files for a given .dll when linking
 - failure to delete them means the VS debugger still holds them locked
 - links: **11, 12, 13, 14**

HCR performance

- snappy compression algorithm - x2-x4 times slower
 - for reference: zlib (c code) to javascript (asm.js) ==> x2 slow down
- calls within a translation unit are direct (the "_actual" version gets called)
- calls between modules => indirection: pointer to function
 - + additional jump from trampoline to actual function
- link time optimization (AKA whole program optimization) cannot help
 - devirtualization techniques are not applicable either
- compactness in memory VS a single binary => instruction cache misses
- /hotpatch for MSVC and **Live++** (which are faster):
 - not going through function pointers
 - by default there are no jumps in the function preamble (padding)
- slowdown depends a lot on the type/scale of software - x2 to x5...

HCR performance

possible optimizations:

- write more "inline" procs
 - their body is emitted wherever used => skip indirections
- pragmas for excluding files (extension of the first point in this list)
 - register the module procs but no indirections between them
- relocate all code from loaded binaries close in memory?
- PLOT TWIST!
 - debug builds are currently affected a lot less (<x2 slowdown)
 - HCR is mainly for development => probably debug builds

HCR TODO

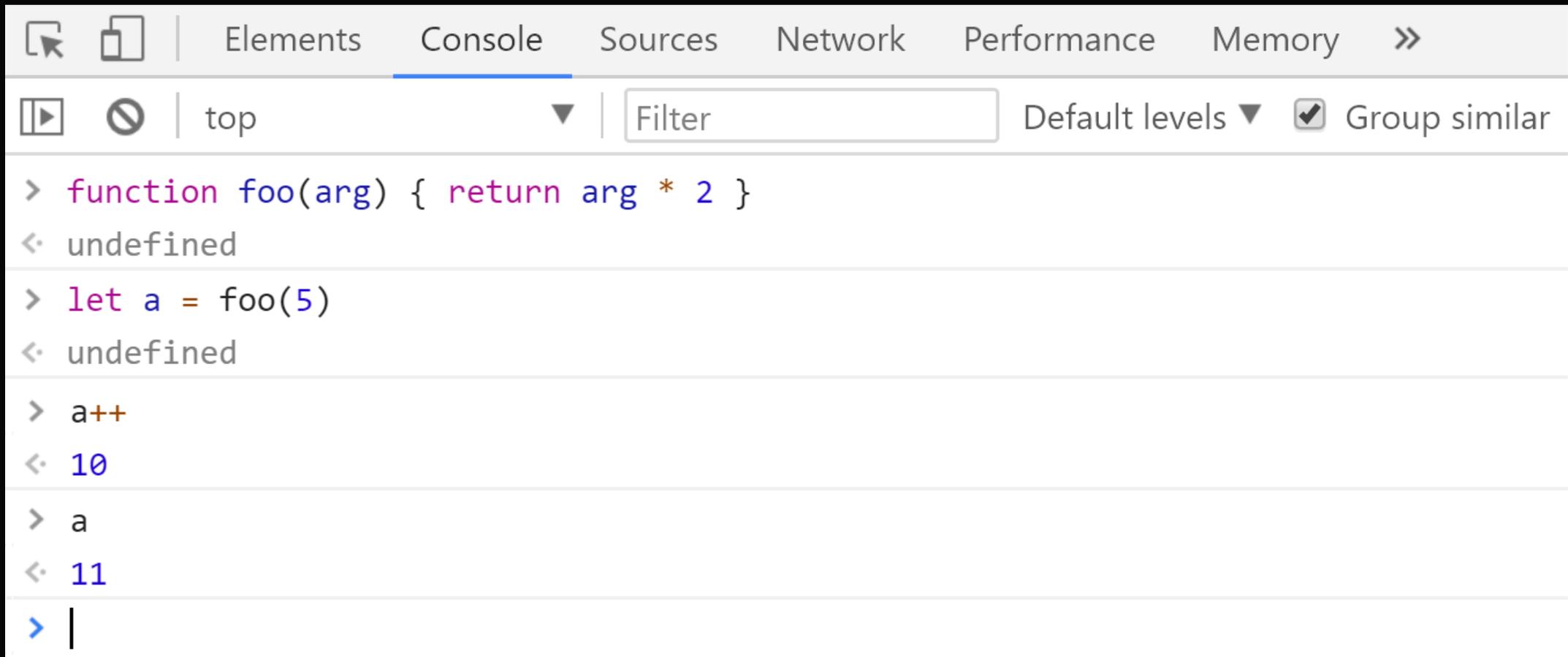
- Nim stdlib has trouble compiling with the GC as a separate SO
 - "-d:useNimRtl" needs to be enabled for all compiler tests
 - currently no real-world project can be built with HCR
- detecting type changes
 - error when detected
 - OR ability for users to handle it (migrate data)
- check if "reload" would affect functions from the current call stack
- expose state for outside manipulation with interactive speeds
 - imagine a slider in the IDE for a variable or a color picker widget
- performance & bug fixes

HCR Implementation choice

- pros
 - any modern (desktop) OS supports dynamic libraries
 - works with any C/C++ compiler
 - near-native speeds
 - final binaries are debuggable
 - a REPL is easily built on top of this
 - (arguably) less complex than using LLVM / JIT / whatever
 - changes are isolated (only the C backend which is a few files)
 - program can be changed in (almost) any way
 - novel approach - someone had to try it
- cons
 - not as optimal as the /hotpatch for MSVC or **Live++**
 - (arguably) more complex than using LLVM / JIT / whatever
 - not sure how NLVM (Nim on top of LLVM) will support HCR

REPL - Read Eval Print Loop

- interpreted languages have it (JavaScript, Python, etc.)
- consoles/shells - cmd.exe, bash
- can iteratively append/execute code (definitions, side effects, etc.)
- education, scientific community, rapid prototyping of any kind



The screenshot shows a web browser's developer console with the 'Console' tab selected. The console displays a series of JavaScript commands and their outputs:

```
> function foo(arg) { return arg * 2 }  
< undefined  
> let a = foo(5)  
< undefined  
> a++  
< 10  
> a  
< 11  
> |
```

REPL/Nim quote

Nim is the language I have always thought was a brilliant idea that I never get to use. It's a shame.

Nim is to C/C++ as CoffeeScript is to JavaScript. A highly extensible template language atop a portable language with libraries for practically everything. So why haven't I hopped on the bandwagon? Outside of C++, C, and Fortran - the only way I have ever learned a new language is through using a REPL. How much of Python's and MATLAB's (and maybe even Julia's) success is due to having a brilliant REPL?

I am not complaining, and I do not have any free time to fix it. But man... if Nim just had a killer REPL that allowed me to slowly learn the language properly while not being blocked from my daily work... it would be just killer!

cjhanks on Apr 18, 2017

<https://news.ycombinator.com/item?id=14143521>

REPL on top of HCR

Talk abstract was a lie! didn't get to implementing it in time...

2 files:

- main module
 - has the main loop
 - handles code submissions
- imported file
 - gets modified based on submissions
 - rebuilt + reloaded

should be well below half a second

REPL on top of HCR

you submit this:

```
import tables  
  
var a = {1: "one", 2: "two"}.toTable  
  
echo a
```

and it gets translated to this:

```
import hotcodere reloading # for the before/after handlers  
  
import tables  
  
var a = {1: "one", 2: "two"}.toTable  
  
afterCodeReload:  
  echo a
```

REPL on top of HCR

later you append:

```
let b = a  
  
echo b
```

and it gets translated to this:

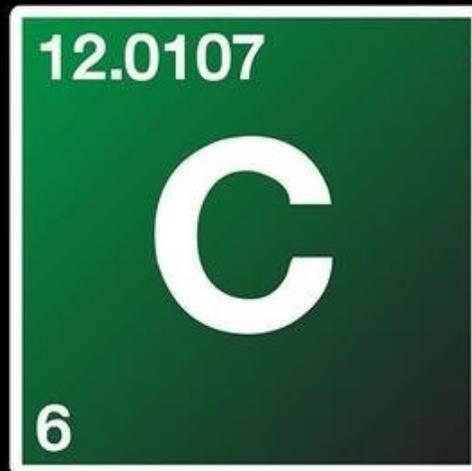
```
import hotcodere reloading # for the before/after handlers  
  
import tables  
  
var a = {1: "one", 2: "two"}.toTable  
  
let b = a # the new code  
  
# only the new side effects are still present  
afterCodeReload:  
  echo b
```

Jupyter kernel

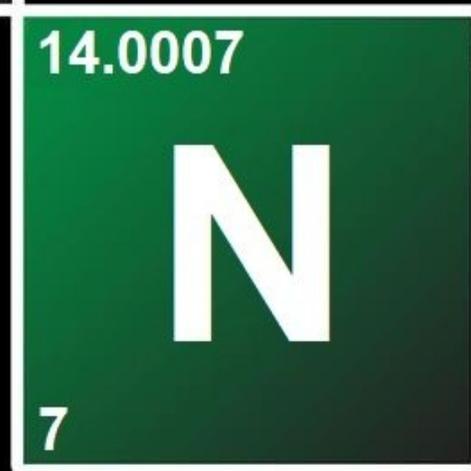
- yesterday on ACCU: Interactive C++ : Meet Jupyter / Cling - The data scientist's geeky younger sibling - by Neil Horlock
- A Jupyter Notebook is an interactive document - a collaborative platform for prototyping, experimentation and analysis
- Mix and share: code, text, data, computation and visualization
- "Notebooks are the most popular tool for working with data at Netflix."
- Nim REPL => Nim Jupyter kernel

The road ahead for Nim

- version 1.0 - promise of stability
- compiler cache for unchanged modules
 - because compilation starts always from the main module
 - of great benefit for HCR/REPL
- more features
- better tooling
- better docs
- taking over the world
- get involved - still in early stages - you can have an impact



hooose



im

Q&A

- <https://nim-lang.org/>
- <https://github.com/nim-lang/Nim>
- FOSDEM 2019: Metaprogramming with Nim

- Slides: https://slides.com/onqtam/nim_hot_code_reloading
- Blog: <http://onqtam.com>
- GitHub: <https://github.com/onqtam>
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