

There's Treasure Everywhere

Prepared for ACCU 2016

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Working Title

How to Write and Deliver a Keynote

What's (in) a Keynote?

- Friend 1: “Just talk about things you like.”
- Friend 2: “Should have no code and no math.”

Things I Like

Things I Like

- Confusing aphorisms

Confusing Aphorisms

When you hear hoofbeats, think
horses, not zebras.

Confusing Aphorisms

When you hear hoofbeats, think
horses, not zebras.

—Not an African Proverb

Things I Like

- Confusing aphorisms

Things I Like

- Confusing aphorisms
- Fast code

Things I Like

- Confusing aphorisms
- Fast code
- C++ Concepts

Toto Washlet



The Toto Company

- Offers 8 washlet models

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- Of these, 7 have a remote control

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- *A remote control*

Washlets with Remotes

- Combine two nice things
 - Complex and interesting
 - Modern and cool
-
- Unlikely to be super useful in practice

“Why Concepts didn’t make C++17”

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“Why Concepts didn’t make C++17”

- Not enough time
- “The Concepts TS does not specify any concept definitions”
- “Use of concepts sometimes results in worse error messages”

Upcoming

“Why Concepts did
make C++20”

Sentinels

Let's talk about linear find

- Nothing out of the ordinary:

```
Range find(Range, E)(Range r, E e) {  
  for (; !r.empty; r.popFront) {  
    if (r.front == e) break;  
  }  
  return r;  
}
```

- Implemented many times in many languages

To index or not to index

- Indexed access and pointer arithmetic changed relative speed over years
- Today: indexed access has a slight edge
 - Less aliasing
 - CPU has sophisticated indexed addressing

Linear find using random access

```
Range find(Range, E)(Range r, E e) {  
    size_t i = 0;  
    for (; i != r.length; ++i) {  
        if (r[i] == e) break;  
    }  
    return r[i .. $];  
}
```

- Task: make it faster

Linear find with sentinel

```
Range find(Range, E)(Range r, E e) {
    auto c = r[$ - 1];
    r[$ - 1] = e;
    size_t i = 0;
    {
        scope(exit) r[$ - 1] = c;
        for (;;) ++i
            if (r[i] == e) break;
    }
    if (i + 1 == r.length && c != e)
        ++i;
    return r[i .. $];
}
```


But... but...

- That doesn't work for all types!
- $r[\$ - 1] = c$ may throw
- Also, `find` itself doesn't work for all types
 - Need to define only when we should
- Need to choose “right” implementation automatically
 - With sentinel for certain types
 - Conservative for the others

Linear find with sentinel, take 2

```
Range find(Range, E)(Range r, E e)
if (isInputRange!Range &&
    is(typeof(r.front) == e) == bool)
) {
    ...
}
```

- `is(typeof(e) == T)` Does expression `e` “work” and has type `T`?
- `e` must *look* like an expression syntactically
- Boolean returned, compilation not stopped
 - SFINAE on steroids
- Still, how to choose with/without sentinel?

static if to the rescue

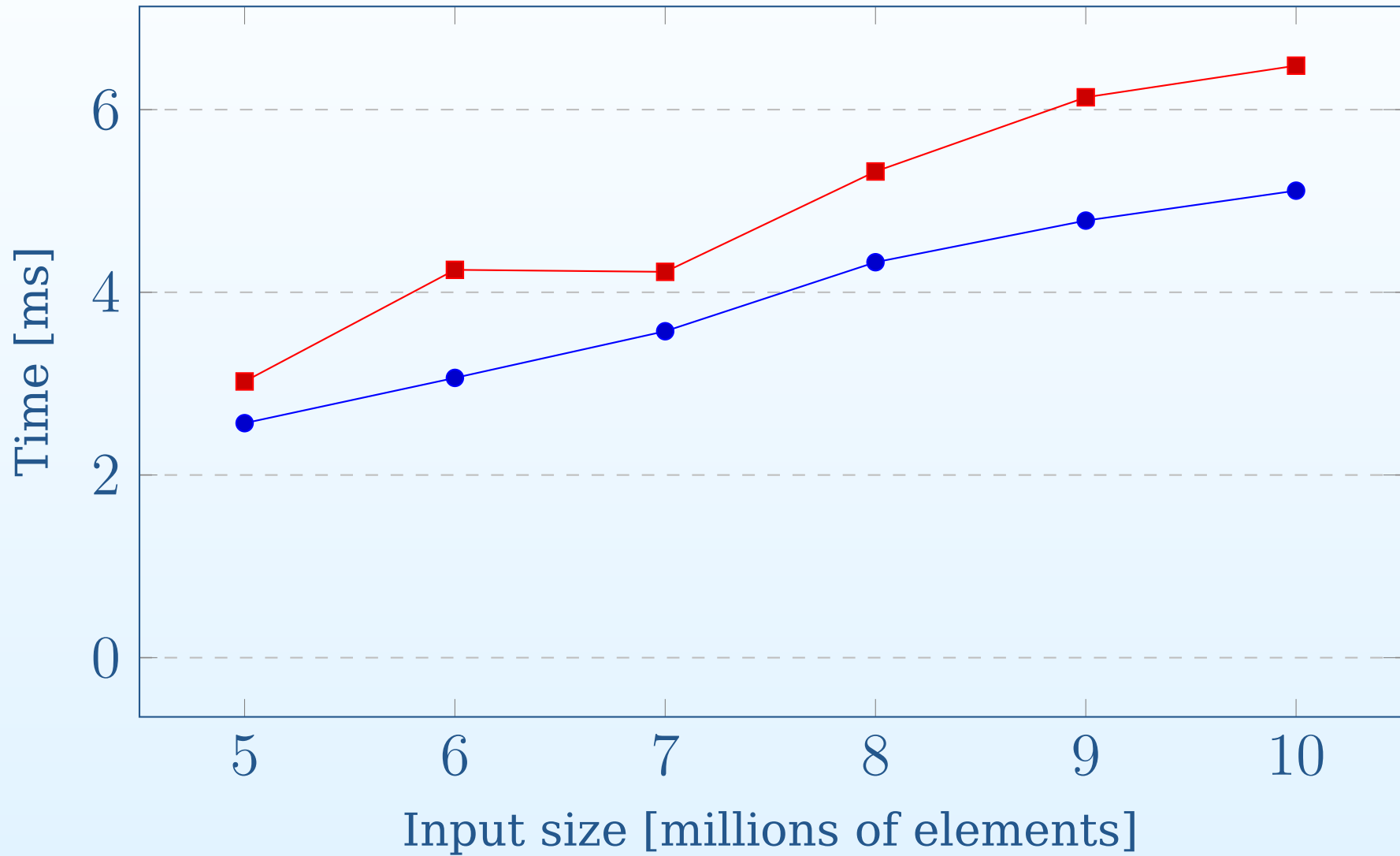
```
Range find(R, E)(R r, E e)
if (isInputRange!R && is(typeof(r.front == e) : bool)) {
    static if (isRandomAccessRange!R && hasSlicing!R) {
        static if (is(typeof(
            () nothrow { r[0] = r[0]; }
        ))) {
            ... sentinel implementation ...
        } else {
            ... indexed implementation ...
        }
    } else {
        ... conservative implementation ...
    }
}
```

Please note

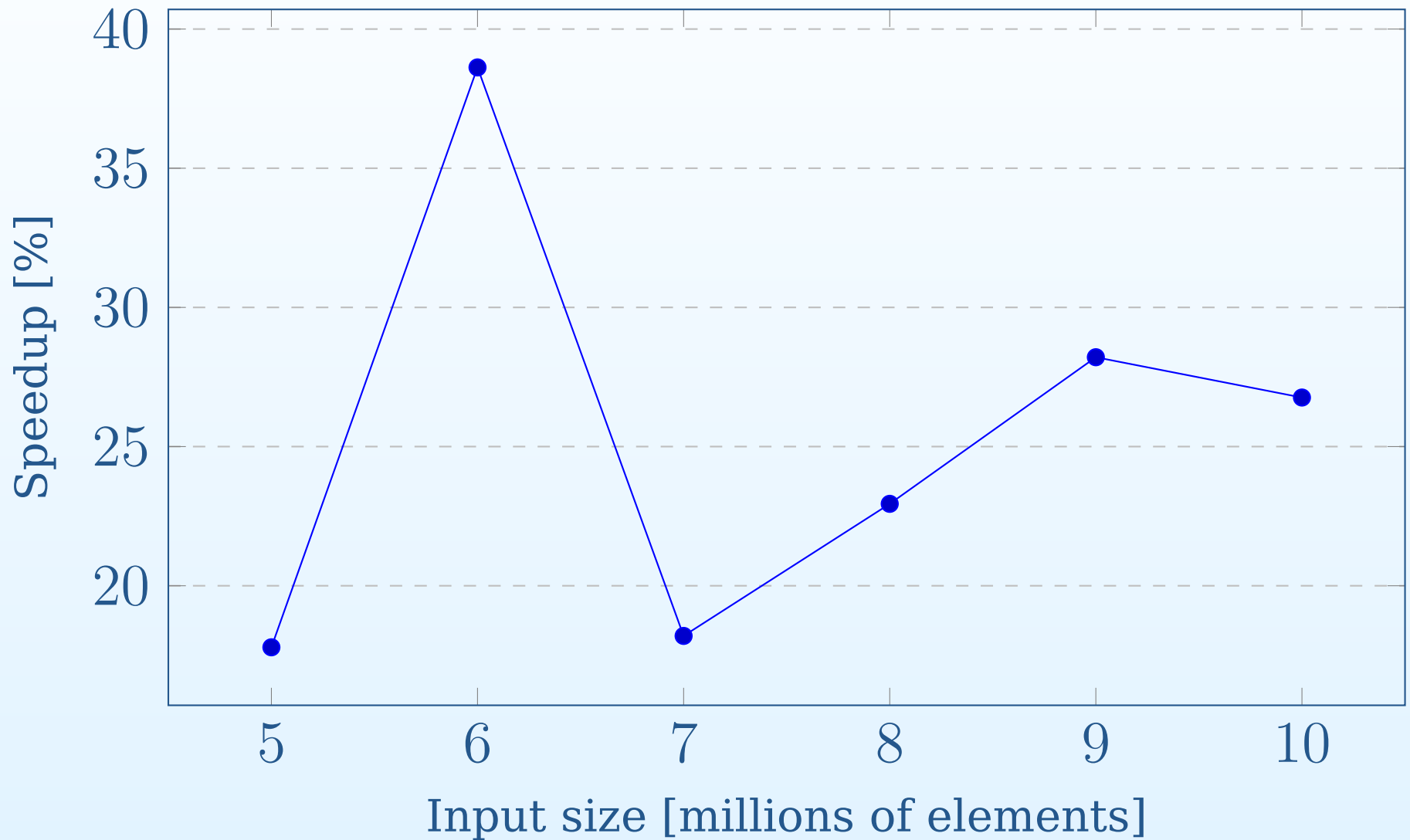
- Sentinel technique old as dust
- Yet not applicable directly to generic code
- Must use proper introspection techniques
- `is` and `static if` widely applicable
 - Mold *structure* depending on capabilities
 - `stdlib`: `static if` every 88 lines as `wc` counts

- Does the optimization help?

Find with Sentinel



Relative Speedup



Sentinels

- Widely applicable technique
- Let's use them for something more complex
- How about Tony Hoare's partition?
 - Workhorse of sorting and the selection problem

Baseline (1/2)

```
size_t pivotPartition(Range)(Range r, size_t pivot) {
    r.swapAt(pivot, 0);
    size_t lo = 1, hi = r.length - 1;
loop: for (;;) lo++, hi-- {
    for (;;) ++lo {
        if (lo > hi) break loop;
        if (r[lo] >= r[0]) break;
    }
    // found the left bound: r[lo] >= r[0]
```


Baseline (2/2)

```
// found the left bound: r[lo] >= r[0]
for (;;) --hi {
    if (lo >= hi) break loop;
    if (r[0] >= r[hi]) break;
}
// found the right bound: r[hi] <= r[0]
// swap them & make progress
r.swapAt(lo, hi);
}
--lo;
r.swapAt(lo, 0);
return lo;
}
```

Make It Faster

Accelerating pivotPartition

- Currently: find from left, find from right, swap, make progress
- Done when indexes meet
- Many index checks interspersed with actual work
- Can't plant a sentinel in the "middle"—final pivot position unknown

Key Idea

- Plant sentinels at both ends
- Create a “vacancy” in the range
- Break swapAt into two assignments
- An assignment fills a vacancy and leaves another one
 Half swap!
- Work becomes matter of moving the vacancy around
- At the end fill back the vacancy

Implementation (1/3): Setup

```
size_t pivotPartition(Range)(Range r, size_t pivot) {  
    if (r.length <= 1) return 0;  
    // Put pivot at the front  
    r.swapAt(pivot, 0);  
    auto p = r[0];  
    // Plant the pivot in the end as well as a sentinel  
    auto save = r[$ - 1];  
    r[$ - 1] = p; // r[$ - 1] is now vacant
```

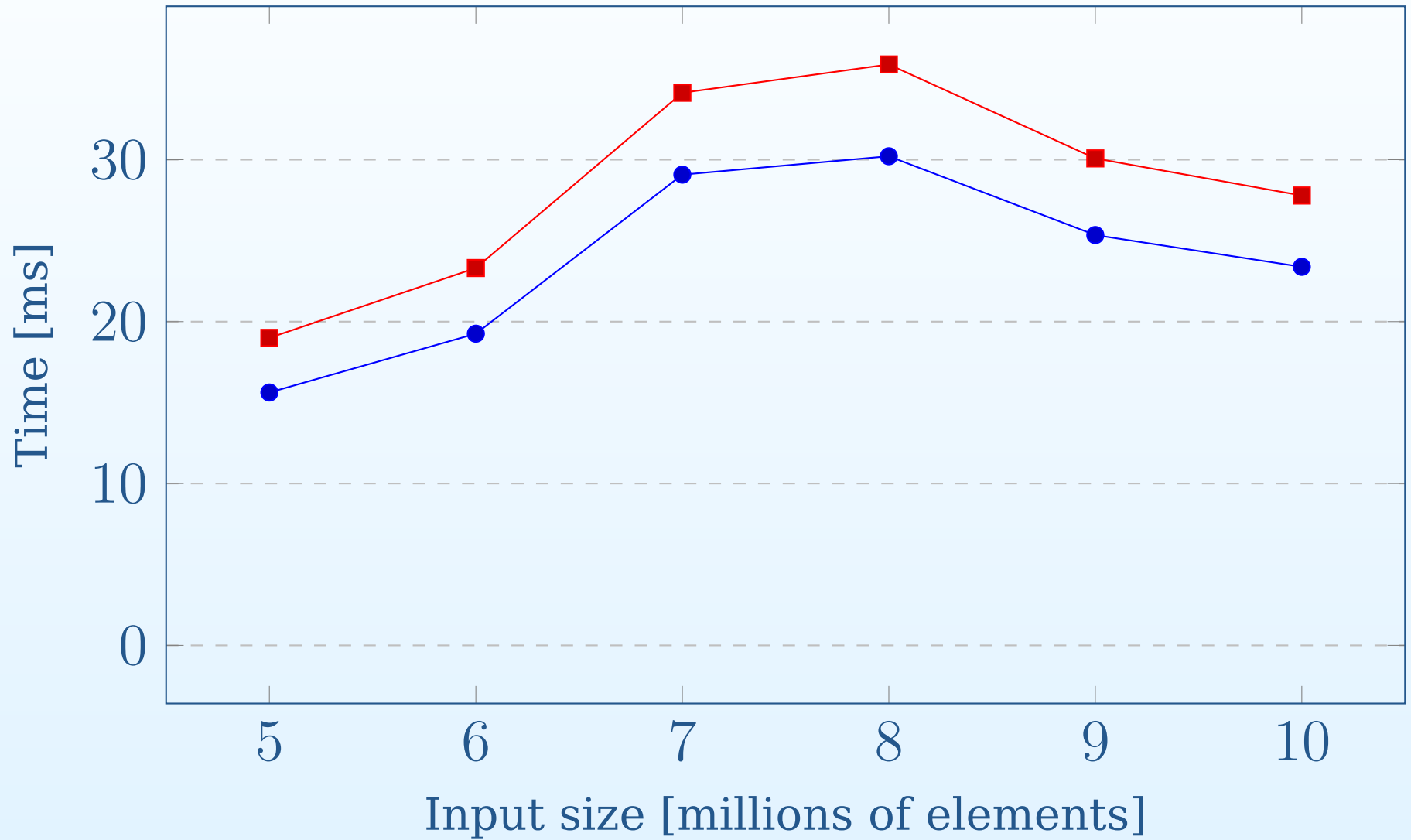
Implementation (2/3): Core

```
size_t lo = 0, hi = r.length - 1;
for (;;) {
    do ++lo; while (r[lo] < p);
    r[hi] = r[lo];
    do --hi; while (p < r[hi]);
    if (lo >= hi) break;
    r[lo] = r[hi];
}
```

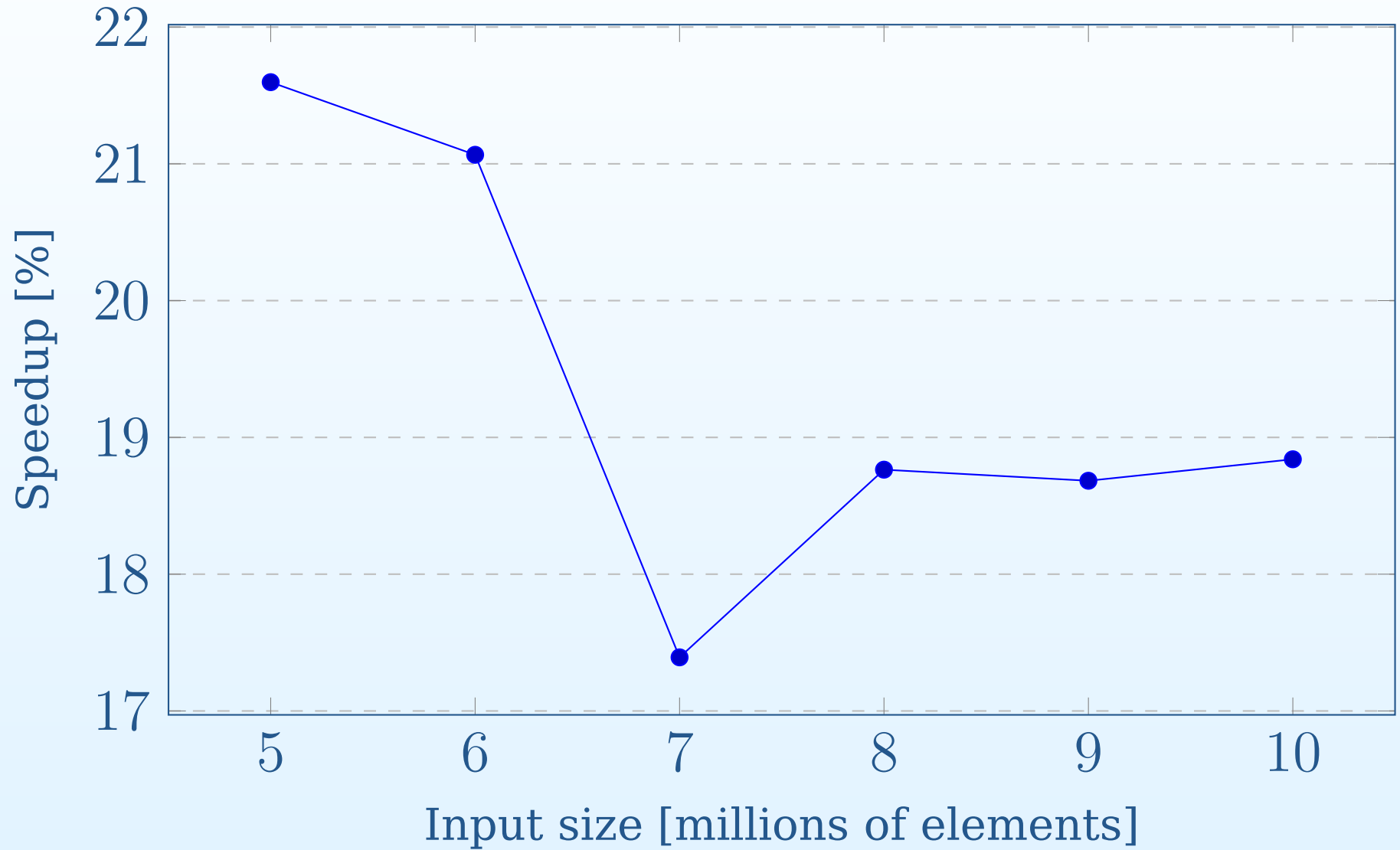
Implementation (3/3): Fixup

```
assert(lo - hi <= 2);
assert(p >= r[hi]);
if (lo == hi + 2) {
    assert(r[hi + 1] >= p);
    r[lo] = r[hi + 1];
    --lo;
}
r[lo] = save;
if (p < save) --lo;
assert(p >= r[lo]);
r.swapAt(0, lo);
return lo;
}
```

Partition



Partition Speedup



Credit

- Thought I invented this!
- Nope: “Engineering of a Quicksort Partitioning Algorithm”, Abhyankar & Ingle, Journal of Global Research in Computer Science, Feb 2011
 - “Vacancy” idea, cannot choose any pivot

More Sentinels

- These sentinel-based optimizations are not cherry picked!
- Dot product
- Set intersection
- Merge sorted lists
- Lexing and parsing

There's Treasure Everywhere

- You just have to find it

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- (by using introspection)

Fastware

- A cornucopia of optimization techniques
 - Benchmarking Techniques
 - Strength Reduction
 - Cache Friendliness
 - Indirect Write Elision
 - Memoization and Obliviation
 - Hoisting and Lowering
 - ... and many more!

Destroy!